

# Simplex Geometry of Graphs

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*in collaboration with Karel Devriendt*

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Google matrix: fundamentals, applications and beyond (GOMAX)  
IHES, October 15-18, 2018



## Outline

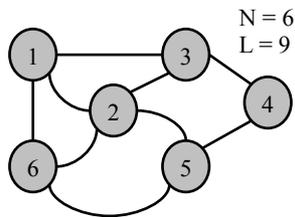


Background:  
Electrical matrix equations

Geometry of a graph



## Adjacency matrix A



$$A_{N \times N} = \begin{bmatrix} 0 & 1 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 & 0 \end{bmatrix}$$

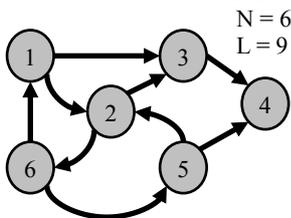
For an undirected graph:  $A = A^T$  is symmetric

Number of neighbors of node  $i$  is the degree:  $d_i = \sum_{k=1}^N a_{ik}$

if there is a link between node  $i$  and  $j$ , then  $a_{ij} = 1$   
else  $a_{ij} = 0$



## Incidence matrix B



- Label links (e.g.:  $l_1 = (1,2)$ ,  $l_2 = (1,3)$ ,  $l_3 = (1,6)$ ,  $l_4 = (2,3)$ ,  $l_5 = (2,5)$ ,  $l_6 = (2,6)$ ,  $l_7 = (3,4)$ ,  $l_8 = (4,5)$ ,  $l_9 = (5,6)$ )
- Col  $k$  for link  $l_k = (i,j)$  is zero, except:  
source node  $i = 1 \rightarrow b_{ik} = 1$   
destination node  $j = -1 \rightarrow b_{jk} = -1$

$$B_{N \times L} = \begin{bmatrix} 1 & 1 & -1 & 0 & 0 & 0 & 0 & 0 & 0 \\ -1 & 0 & 0 & 1 & -1 & 1 & 0 & 0 & 0 \\ 0 & -1 & 0 & -1 & 0 & 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 0 & -1 & -1 & 0 \\ 0 & 0 & 0 & 0 & 1 & 0 & 0 & 1 & -1 \\ 0 & 0 & 1 & 0 & 0 & -1 & 0 & 0 & 1 \end{bmatrix}$$

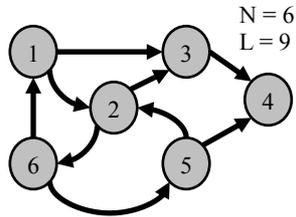
Col sum B is zero:  $u^T B = 0$

where the all-one vector  $u = (1,1,\dots,1)$

B specifies the directions of links



## Laplacian matrix Q



$$Q_{N \times N} = \begin{bmatrix} 3 & -1 & -1 & 0 & 0 & -1 \\ -1 & 4 & -1 & 0 & -1 & -1 \\ -1 & -1 & 3 & -1 & 0 & 0 \\ 0 & 0 & -1 & 2 & -1 & 0 \\ 0 & -1 & 0 & -1 & 3 & -1 \\ -1 & -1 & 0 & 0 & -1 & 3 \end{bmatrix}$$

$$Q = BB^T = \Delta - A$$

$$\Delta = \text{diag}(d_1 \quad d_2 \quad \dots \quad d_N)$$

Since  $BB^T$  is symmetric, so are  $A$  and  $Q$ . Although  $B$  specifies directions,  $A$  and  $Q$  lost this info here.

Basic property:  $Qu = 0$

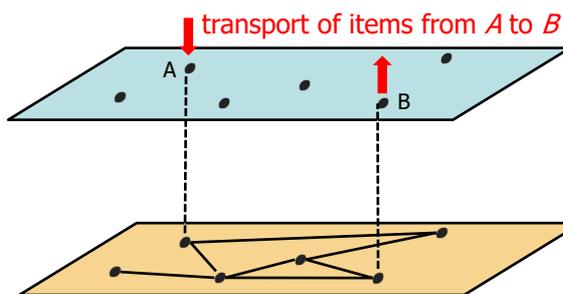
$u$  is an eigenvector of  $Q$   
Belonging to eigenvalue  $\mu = 0$

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$$Qu = BB^T u = 0 \quad \text{because} \quad 0 = u^T B = B^T u$$



## Network: service(s) + topology



### Service (function)

software, algorithms

### Topology (graph)

hardware, structure

### Service and topology

- own specifications
- both are, generally, time-variant
- service is often designed independently of the topology
- often more than 1 service on a same topology

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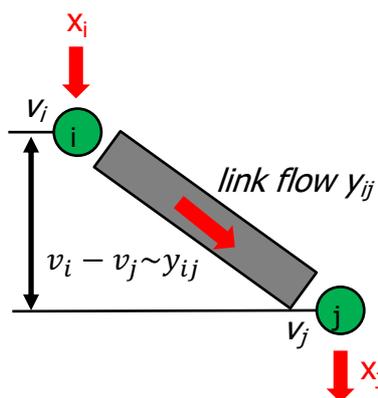
## Function of network

- Usually, the function of a network is related to the *transport of items over its underlying graph*
- In man-made infrastructures: two major types of transport
  - Item is a **flow** (e.g. electrical current, water, gas,...)
  - Item is a **packet** (e.g. IP packet, car, container, postal letter,...)
- **Flow equations (physical laws)** determine transport (Maxwell equations (Kirchhoff & Ohm), hydrodynamics, Navier-Stokes equation (turbulent, laminar flow equations, etc.))
- **Protocols** determine transport of packets (IP protocols and IETF standards, car traffic rules, etc.)

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## Linear dynamics on networks

Linear dynamic process: "proportional to" ( $\sim$ ) graph of network



### Examples:

- water (or gas) flow  $\sim$  pressure
- displacement (in spring)  $\sim$  force
- heat flow  $\sim$  temperature
- **electrical current  $\sim$  voltage**

$$x = Q \cdot v$$

injected nodal current vector	=	weighted Laplacian of the graph	.	nodal potential vector
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## Pseudoinverse of the Laplacian (review)

The inverse of the current-voltage relation  $\mathbf{x} = \mathbf{Q}\mathbf{v}$   
 is the voltage-current relation  $\mathbf{v} = \mathbf{Q}^\dagger \mathbf{x}$   
 subject to  $\mathbf{u}^T \mathbf{x} = 0$  and  $\mathbf{u}^T \mathbf{v} = 0$

The spectral decomposition

$$\tilde{\mathbf{Q}} = \sum_{k=1}^{N-1} \tilde{\mu}_k \mathbf{z}_k \mathbf{z}_k^T$$

allows us to compute the pseudoinverse (or Moore-Penrose inverse)

$$\mathbf{Q}^\dagger = \sum_{k=1}^{N-1} \frac{1}{\tilde{\mu}_k} \mathbf{z}_k \mathbf{z}_k^T$$

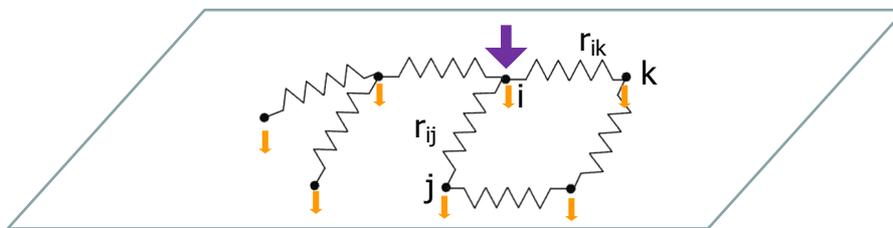
The effective resistance  $N \times N$  matrix is  $\tilde{\Omega} = \mathbf{u} \zeta^T + \zeta \mathbf{u}^T - 2\mathbf{Q}^\dagger$ ,  
 where the  $N \times 1$  vector  $\zeta = (Q_{11}^\dagger, Q_{22}^\dagger, \dots, Q_{NN}^\dagger)$

An interesting graph metric is the effective graph resistance

$$R_G = \mathbf{N} \mathbf{u}^T \zeta = N \text{trace}(\mathbf{Q}^\dagger) = N \sum_{k=1}^{N-1} \frac{1}{\mu_k}$$

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P. Van Mieghem, K. Devriendt and H. Cetinay, 2017, "Pseudo-inverse of the Laplacian and best spreader node in a network", Physical Review E, vol. 96, No. 3, p 032311.



Inverses:  $\mathbf{x} = \mathbf{Q}\mathbf{v} \leftrightarrow \mathbf{v} = \mathbf{Q}^\dagger \mathbf{x}$  with voltage reference  $\mathbf{u}^T \mathbf{v} = 0$

$\mathbf{Q}^\dagger$  : pseudoinverse of the weighted Laplacian obeying  $\mathbf{Q}\mathbf{Q}^\dagger = \mathbf{Q}^\dagger \mathbf{Q} = \mathbf{I} - \frac{1}{N} \mathbf{J}$   
 $\mathbf{J} = \mathbf{u}\mathbf{u}^T$  : all-one matrix  $\mathbf{u}$  : all-one vector

Unit current injected in node  $i$   
 $\mathbf{x} = \mathbf{e}_i - \frac{1}{N} \mathbf{u}$



nodal potential of  $i$   
 $v_i = Q_{ii}^\dagger$

The best spreader is node  $k$  with  $Q_{kk}^\dagger \leq Q_{ii}^\dagger$  for  $1 \leq i \leq N$



## Outline



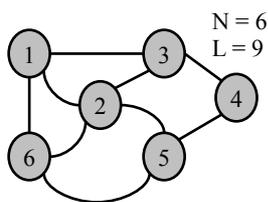
Background:  
Electrical matrix equations

Geometry of a graph



## Three representations of a graph

Topology domain



$$A_{N \times N} = \begin{bmatrix} 0 & 1 & 1 & 0 & 0 & 1 \\ 1 & 0 & 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 & 0 & 0 \\ 0 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 0 & 1 & 0 & 1 \\ 1 & 1 & 0 & 0 & 1 & 0 \end{bmatrix}$$

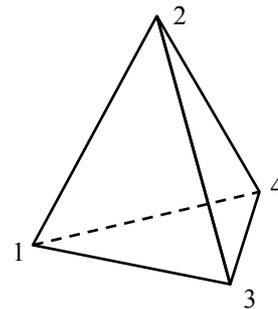
Spectral domain

$$A = A^T = X \Lambda X^T$$

$X_{N \times N}$ : orthogonal  
eigenvector matrix

$\Lambda_{N \times N}$ : diagonal  
eigenvalue matrix

Geometric domain



Undirected graph on  
 $N$  nodes  
= **simplex** in Euclidean  
 $(N-1)$ -dimensional space

Devriendt, K. and P. Van Mieghem, 2018, "The Simplex Geometry of Graphs",  
Delft University of Technology, report20180717.  
(<http://arxiv.org/abs/1807.06475>).



## Miroslav Fiedler (1926-2015)

### MATRICES AND GRAPHS IN GEOMETRY

Miroslav Fiedler



Father of "algebraic connectivity"

His 1972 paper: > 3400 citations

*"This book comprises, in addition to auxiliary material, the research on which I have worked for over 50 years."*

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more information - [www.cambridge.org/9780521461931](http://www.cambridge.org/9780521461931)

appeared in 2011



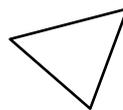
## What is a simplex?



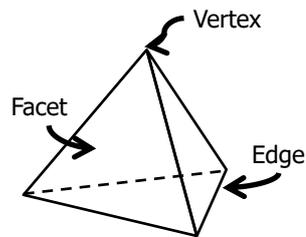
Point



Line Segment



Triangle



Tetrahedron

*Roughly* : a simplex is generalization of a triangle to  $N$  dimensions

*Potential* : Euclidean geometry is the oldest, mathematical theory

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## Spectral decomposition weighted Laplacian (1)

Spectral decomposition:  $Q = ZMZ^T$

where  $M = \text{diag}(\mu_1, \mu_2, \dots, \mu_{N-1}, 0)$ , because  $Q \mathbf{1} = 0$

and the eigenvector matrix  $Z$  obeys  $Z^T Z = Z Z^T = I$  with structure

$$\text{node } \mathbf{Z} = \begin{bmatrix} (z_1)_1 & (z_2)_1 & \cdots & (z_N)_1 \\ (z_1)_2 & (z_2)_2 & \cdots & (z_N)_2 \\ \vdots & \vdots & \ddots & \vdots \\ (z_1)_N & (z_2)_N & \cdots & (z_N)_N \end{bmatrix} = \begin{bmatrix} (z_1)_1 & (z_2)_1 & \cdots & 1/\sqrt{N} \\ (z_1)_2 & (z_2)_2 & \cdots & 1/\sqrt{N} \\ \vdots & \vdots & \ddots & \vdots \\ (z_1)_N & (z_2)_N & \cdots & 1/\sqrt{N} \end{bmatrix}$$

↑  
frequencies  
(eigenvalues)

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$$Q = \sum_{k=1}^{N-1} \mu_k z_k z_k^T$$

TU Delft

## Spectral decomposition weighted Laplacian (2)

Only for a positive semi-definite matrix, it holds that

$$Q = ZMZ^T = (Z\sqrt{M})(Z\sqrt{M})^T$$

The matrix  $S = (Z\sqrt{M})^T$  obeys  $Q = S^T S$  and has rank  $N-1$   
(row  $N = 0$  due to  $\mu_N = 0$ )

$$S = \begin{bmatrix} \sqrt{\mu_1}(z_1)_1 & \sqrt{\mu_1}(z_1)_2 & \cdots & \sqrt{\mu_1}(z_1)_N \\ \sqrt{\mu_2}(z_2)_1 & \sqrt{\mu_2}(z_2)_2 & \cdots & \sqrt{\mu_2}(z_2)_N \\ \vdots & \vdots & \ddots & \vdots \\ \sqrt{\mu_{N-1}}(z_{N-1})_1 & \sqrt{\mu_{N-1}}(z_{N-1})_2 & \cdots & \sqrt{\mu_{N-1}}(z_{N-1})_N \\ 0 & 0 & \cdots & 0 \end{bmatrix}$$

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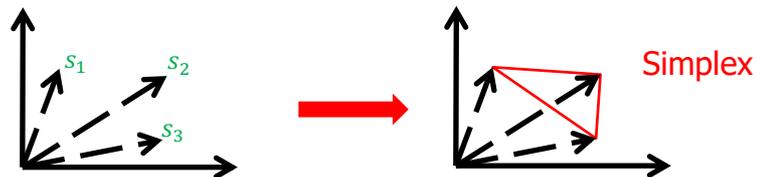
$$Q = \sum_{k=1}^{N-1} \mu_k z_k z_k^T$$

TU Delft

## Geometrical representation of a graph

$$S = \begin{bmatrix} \sqrt{\mu_1}(z_1)_1 & \sqrt{\mu_1}(z_1)_2 & \cdots & \sqrt{\mu_1}(z_1)_N \\ \sqrt{\mu_2}(z_2)_1 & \sqrt{\mu_2}(z_2)_2 & \cdots & \sqrt{\mu_2}(z_2)_N \\ \vdots & \vdots & \ddots & \vdots \\ \sqrt{\mu_{N-1}}(z_{N-1})_1 & \sqrt{\mu_{N-1}}(z_{N-1})_2 & \cdots & \sqrt{\mu_{N-1}}(z_{N-1})_N \\ \hline 0 & 0 & \cdots & 0 \end{bmatrix}$$

The  $i$ -th column vector  $s_i = (\sqrt{\mu_1}(z_1)_i, \sqrt{\mu_2}(z_2)_i, \dots, \sqrt{\mu_{N-1}}(z_{N-1})_i, 0)$  represents a point  $p_i$  in  $(N-1)$ -dim space (because  $S$  has rank  $N-1$ )



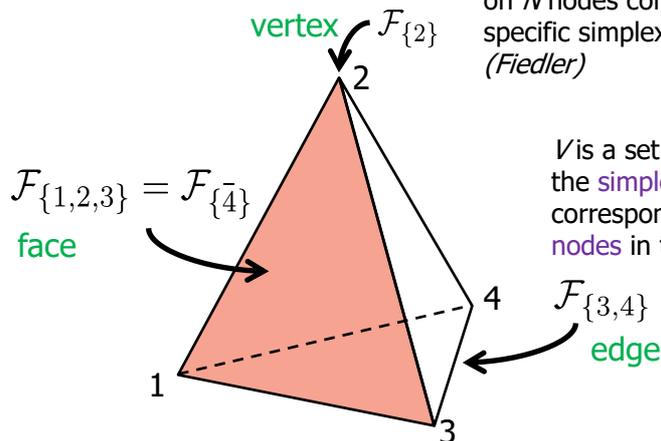
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Simplex geometry: omit zero row,  $S_{N \times N} \rightarrow S_{(N-1) \times N}$



## Faces of a simplex

Each connected, undirected graph on  $N$  nodes corresponds to 1 specific simplex in  $N-1$  dimensions (Fiedler)



$V$  is a set of vertices of the simplex in  $\mathbb{R}^{N-1}$ , corresponding to a set of nodes in the graph  $G$

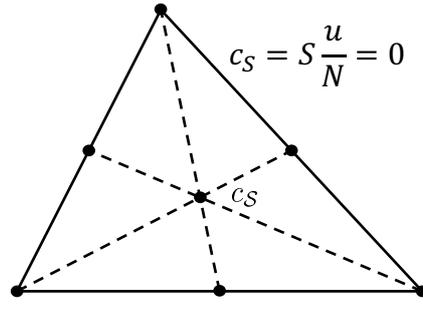
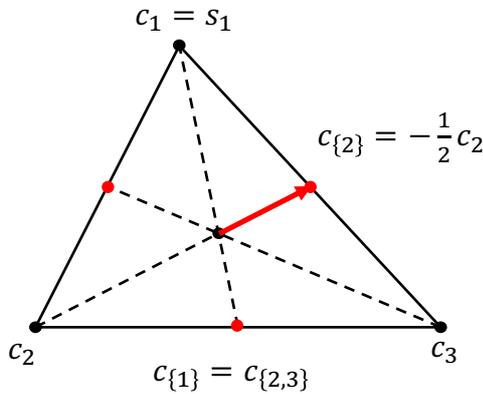
A face  $F_V = \{p \in \mathbb{R}^{N-1} | p = Sx_V \text{ with } (x_V)_i \geq 0 \text{ and } u^T x_V = 1\}$

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The vector  $x_V \in \mathbb{R}^N$  is a **barycentric coordinate** with  $\begin{cases} (x_V)_i \in \mathbb{R} & \text{if } i \in V \\ (x_V)_i = 0 & \text{if } i \notin V \end{cases}$

## Centroids

$c_V = S \frac{u_V}{|V|}$  is the centroid of face  $F_V$  with  $(u_V)_i = 1_{i \in V}$



a centroid of a face is a vector

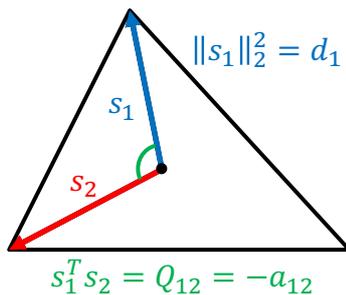
centroid of simplex is origin

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$u_V = u - u_V \rightarrow |V|c_V = S(u - u_V) = -(N - V)c_V$

## Geometric representation of a graph

$$\begin{aligned} \|s_i\|_2^2 &= d_i & \|s_i - s_j\|_2^2 &= (s_i - s_j)^T (s_i - s_j) = s_i^T s_i + s_j^T s_j - 2s_i^T s_j \\ & & &= Q_{ii} + Q_{jj} - 2Q_{ij} \\ & & &= d_i + d_j + 2a_{ij} \text{ for } i \neq j, \text{ else zero} \end{aligned}$$



The matrix with off-diagonal elements  $d_i + d_j + 2a_{ij}$  is a **distance matrix** (if the graph  $G$  is connected)

The geometric graph representation is not unique (node relabeling changes  $Z$ )

$$s_i^T s_j = \sum_{k=1}^{N-1} \sqrt{\mu_k} (z_k)_i \sqrt{\mu_k} (z_k)_j = \sum_{k=1}^{N-1} \mu_k (z_k z_k^T)_{ij} = Q_{ij}$$

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$Q = \sum_{k=1}^{N-1} \mu_k z_k z_k^T$  and  $Q = S^T S$

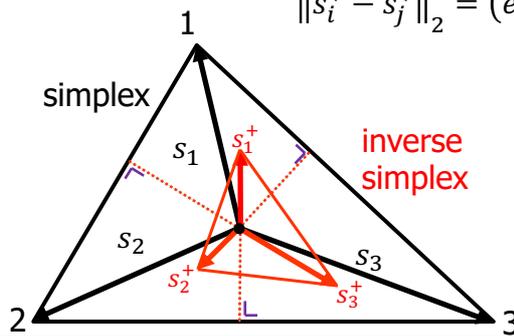
## Geometry of a graph (dual representation)

Spectral decomposition:  $Q^\dagger = ZM^\dagger Z^T = (Z\sqrt{M^\dagger})(Z\sqrt{M^\dagger})^T$

The matrix  $S^\dagger = (Z\sqrt{M^\dagger})^T$  has rank  $N-1$  and  $Q^\dagger = (S^\dagger)^T S^\dagger$

The  $i$ -th column vector  $s_i^\dagger$  obeys

$$\|s_i^\dagger - s_j^\dagger\|_2^2 = (e_i - e_j)^T Q^\dagger (e_i - e_j) = \omega_{ij}$$



From  $S^T S^\dagger = I - \frac{uu^T}{N}$ :

$$\begin{cases} s_j^T s_i^\dagger = -\frac{1}{N} \\ s_i^T s_i^\dagger = 1 - \frac{1}{N} \end{cases}$$

$$\Rightarrow (s_k - s_j)^T s_i^\dagger = 0$$

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$\omega_{ij}$  is the effective resistance between node  $i$  and  $j$



## Volume of simplex and inverse simplex of a graph

Volume of the simplex  $V_G = \frac{N}{(N-1)!} \sqrt{\xi}$

where the number of (weighted) spanning trees  $\xi$  is  $\xi = \frac{1}{N} \prod_{k=1}^{N-1} \mu_k$

Volume of the inverse simplex  $V_G^\dagger = \frac{1}{(N-1)! \sqrt{\xi}}$

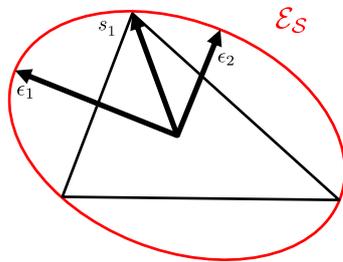
Hence:  $\frac{V_G}{V_G^\dagger} = N\xi = \prod_{k=1}^{N-1} \mu_k$

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K. Menger, "New foundation of Euclidean geometry", American Journal of Mathematics, 53(4):721-745, 1931



## Steiner ellipsoid of simplex



projection  $s_1^T \epsilon_2 = \mu_2(z_2)_1$

semi-axis:  $\|\epsilon_2\| = \sqrt{\frac{N}{N-1}} \mu_2$

volume:

$$V_{\mathcal{E}_S} = \frac{\pi^{N/2}}{\Gamma(N/2+1)} \frac{N^{N/2}}{(N-1)^{N/2}} \sqrt{\prod_{k=1}^N \mu_k}$$

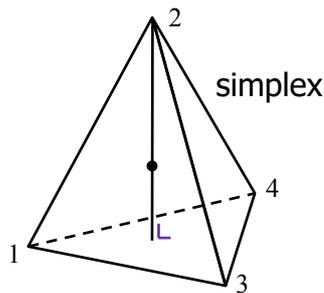
Hence,

$$V_{\mathcal{E}_S}^2 = \frac{(N\pi)^N}{(\Gamma(N/2+1))^2 (N-1)^N} \prod_{k=1}^N \mu_k$$

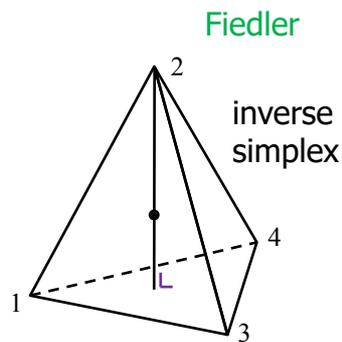
$$V_{N\text{-ellipsoid}} = \frac{\pi^{N/2}}{\Gamma(N/2+1)} \prod_{k=1}^N \epsilon_k$$



## altitude(s) in a simplex



$$\|a_{\{2\}}\|^2 = \frac{1}{Q_{22}^+}$$



$$\|a_2^+\|^2 = \frac{1}{d_2}$$

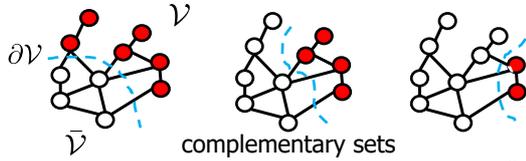
The altitude from a vertex  $s_i^+$  to the complementary face  $F_{\{i\}}^+$  in the inverse simplex (dual graph representation) has a length equal to the inverse degree of node  $i$

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recall that  $Q_{ii}^+ = v_i$  (nodal potential, best spreader)



## Cut size of graph and altitude of simplex



$$\text{Cut size: } |\partial V| = u_V^T Q u_V$$

$$Q = S^T S$$

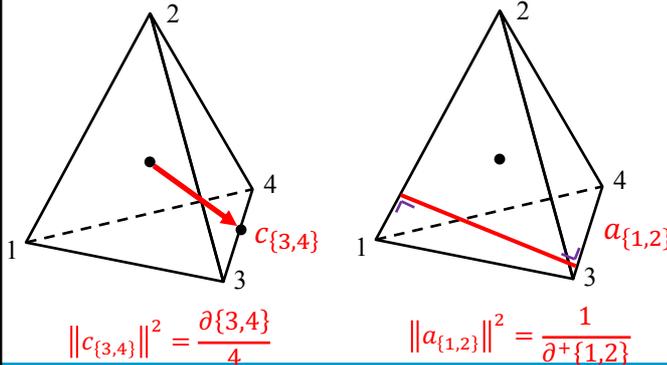
$$\begin{aligned} |\partial V| &= u_V^T S^T S u_V \\ &= \|S u_V\|_2^2 \\ &= \|c_V\|_2^2 V^2 \end{aligned}$$

altitude:

vector from face  $F_V$  to face  $F_V$  and orthogonal to both faces

$$|\partial^+ V| = u_V^T Q^+ u_V$$

$$\|a_V^+\|^2 = \frac{1}{|\partial V|}$$



$$\|c_{\{3,4\}}\|^2 = \frac{\partial\{3,4\}}{4}$$

$$\|a_{\{1,2\}}\|^2 = \frac{1}{\partial^+\{1,2\}}$$

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$c_V = S \frac{u_V}{|V|}$  is the centroid of face  $F_V$  with  $(u_V)_i = 1_{i \in V}$



## Metrics $\sqrt{\omega_{ij}}$ and $\omega_{ij}$

$$\|s_i^\dagger - s_j^\dagger\|_2 = \sqrt{\omega_{ij}}$$

the Euclidean distance between vertices of inverse simplex  $S^\dagger$



vertices of  $S^\dagger$  are an embedding of nodes of the graph  $G$  according to the metric  $\sqrt{\omega_{ij}}$  (a.o. obeying the triangle inequality)

$$\text{Also } \|s_i^\dagger - s_j^\dagger\|_2^2 = \omega_{ij} \text{ is a metric}$$



Inverse simplex  $S^\dagger$  of the graph  $G$  with positive link weights is hyperacute

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Unpublished recent work with K. Devriendt



## Generalization metrics

$Q^\dagger$  is the Gram matrix of a hyperacute simplex  $S^+$   $\rightarrow$  determines a metric



$m_{ij}^{(f)} = (e_i - e_j)^T (f(Q))^\dagger (e_i - e_j)$  is a metric when  $f(Q)$  is a Laplace matrix

“Statistical physics” metrics on a graph

$$m_{ij}^{(stat(g))} = (e_i - e_j)^T (e^{(Q+pQ_K)t} - I + gQ_K)^\dagger (e_i - e_j) = \sum_{k=1}^{N-1} \frac{(z_k)_i - (z_k)_j}{e^{(\mu_k + pN)t + g}}$$

with  $t = \frac{1}{k_B T}$  and  $pN = -E_f$  (chemical potential or Fermi energy)

$g = -1 \rightarrow$  Bose-Einstein

$g = 0 \rightarrow$  Maxwell-Boltzmann

$g = 1 \rightarrow$  Fermi-Dirac

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Unpublished recent work with K. Devriendt



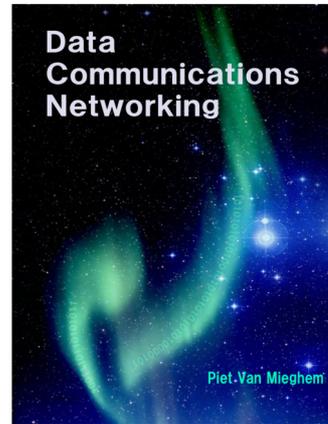
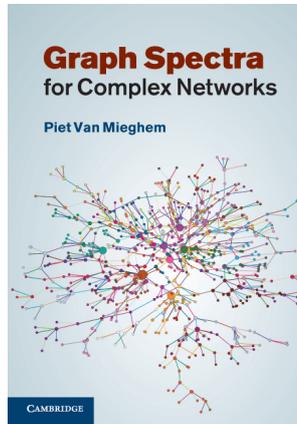
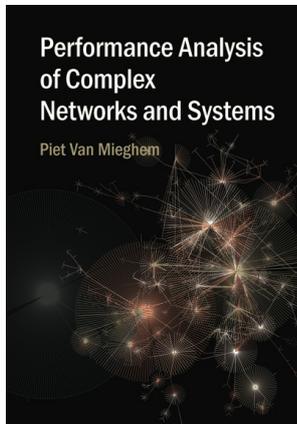
## Summary

- Linearity between process and graph naturally leads to the weighted Laplacian  $Q$  and its pseudoinverse  $Q^\dagger$
- Spectral decomposition of the weighted Laplacian  $Q$  and its pseudoinverse  $Q^\dagger$  provides an  $N-1$  dimensional simplex representation of each graph,
  - allowing computations in the  $N-1$  dim. Euclidean space (in which a distance/norm is defined)
  - geometry for (undirected) graphs
- **Open:** “Which network problems are best solved in the simplex representation?”

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## Books



Articles: <http://www.nas.ewi.tudelft.nl>

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# Thank You

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